

# Monoidal Theories and Graph Rewriting in Quantum Computing

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- ▶ Most importantly, we have a graphical notation
- ▶ Applications to quantum states and processes motivates the study of graphical, monoidal theories in categories like  $\text{FdHilb}$
- ▶ To do graphical proofs (esp. automatic ones), we need the correct notion of graph rewriting, as well as sound, efficient algorithms

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- ▶ Mixed state quantum mechanics has generalisations of the above. We won't talk about that.

# Entanglement and the Tensor

- ▶ For our purposes, take  $\otimes$  to be the Kronecker product:

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- ▶ For Hilbert spaces  $\mathcal{H}_1$  and  $\mathcal{H}_2$ , we can construct  $\mathcal{H}_1 \otimes \mathcal{H}_2 = \text{span} \{v \otimes u : v \in \mathcal{H}_1, u \in \mathcal{H}_2\}$ .
  - ▶  $\dim(\mathcal{H}_1 \otimes \mathcal{H}_2) = \dim \mathcal{H}_1 \cdot \dim \mathcal{H}_2$



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- ▶ The Hilbert space  $\mathcal{Q} := \mathbb{C}^2$  is called the space of *qubits*.
- ▶ We write the standard basis of  $\mathcal{Q}$  in “ket” notation, as  $|0\rangle, |1\rangle$ . Also,  $|ij\rangle$  is shorthand for  $|i\rangle \otimes |j\rangle$ .

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- ▶  $(\text{FdHilb}, \otimes, \mathbb{C}, (-)^{\dagger})$  is a dagger symmetric monoidal category

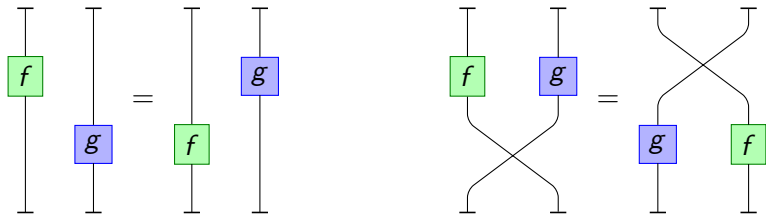
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- ▶ Other examples include  $\text{Rel}$ ,  $\text{nCob}$ ,  $\text{Span}(\mathcal{C})$ ,  $\text{Csp}(\mathcal{C})$ , group representation categories, projective spaces, ...

# Graphical Representation

$$f \otimes g := \begin{array}{c} \top \\ | \\ \boxed{f} \\ | \\ \text{---} \end{array} \quad \begin{array}{c} \top \\ | \\ \boxed{g} \\ | \\ \text{---} \end{array} \qquad g \circ f := \begin{array}{c} \top \\ | \\ \boxed{f} \\ | \\ \boxed{g} \\ | \\ \text{---} \end{array}$$

- ▶ We can express the bifactoriality of  $\otimes$  and the naturality of  $\sigma$  as follows:



# Quantum Circuits

- ▶ Any †-SMC provides a notion of *unitarity* ( $U^\dagger U = 1$  and  $UU^\dagger = 1$ )
- ▶ We can think of unitary maps as the quantum analogy to reversible logic gates
- ▶ Lets look at one motivating example:

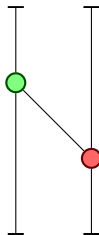
$$\text{CNOT} :: |00\rangle \rightarrow |00\rangle, |01\rangle \rightarrow |01\rangle, |10\rangle \rightarrow |11\rangle, |11\rangle \rightarrow |10\rangle$$

- ▶ Usually drawn like this:



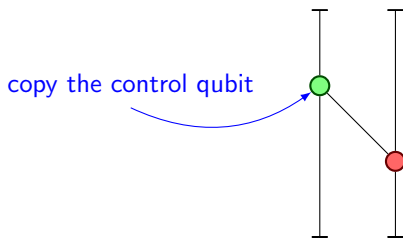
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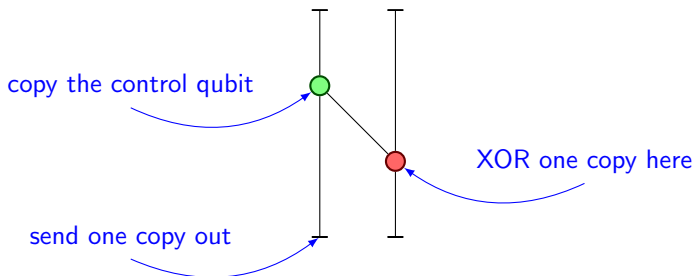






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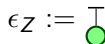
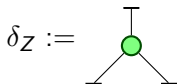


# Classical Structures

- ▶ A chosen basis is like some classical data embedded in the system.
- ▶ What can we do with classical data?
  - ▶ Copy and delete!

$$\delta_Z : \mathcal{Q} \rightarrow \mathcal{Q} \otimes \mathcal{Q} :: |i\rangle \mapsto |ii\rangle \qquad \epsilon_Z : \mathcal{Q} \rightarrow \mathbb{C} :: |i\rangle \mapsto 1$$

- ▶ Graphically:



# Classical Structures

- ▶  $(-)^{\dagger}$  flips everything upside-down:

$$\delta_Z^{\dagger} := \begin{array}{c} \diagup \quad \diagdown \\ \text{●} \\ \text{I} \end{array} \quad \epsilon_Z^{\dagger} := \begin{array}{c} \text{●} \\ \text{I} \end{array}$$

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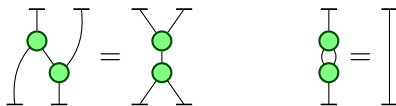
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- ▶ Likewise,  $(\delta_Z, \epsilon_Z)$  is an *internal cocommutative comonoid*

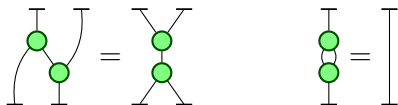
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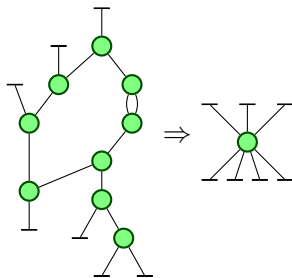


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- ▶ Terms of these things are uniquely determined by number of inputs and outputs. As a result, we write connected graphs:





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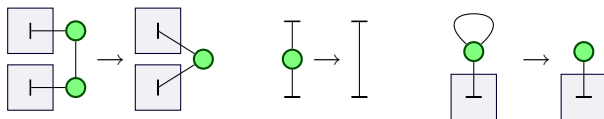
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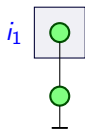


## More on !-boxes

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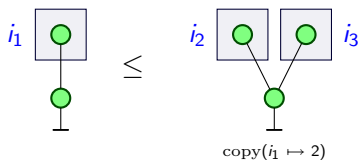
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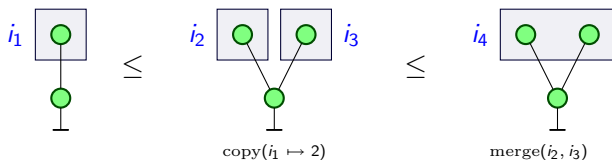
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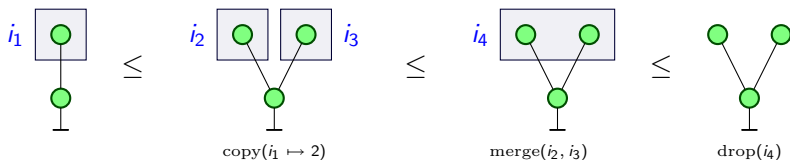
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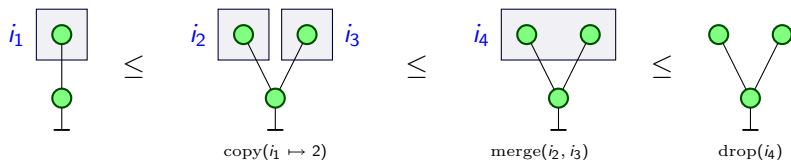
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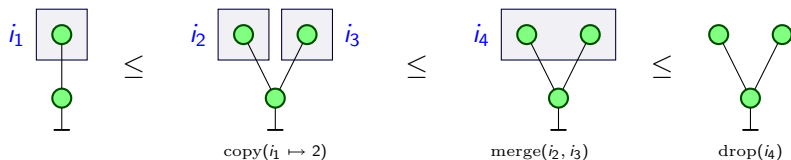
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- ▶ Naive rewriting algorithm is to expand all combinations of !-boxes, bounded by the size of the target
- ▶ Better algorithm expands lazily during matching

## Another colour

- ▶ Let  $(\delta_X, \epsilon_X)$  copy and delete this basis:

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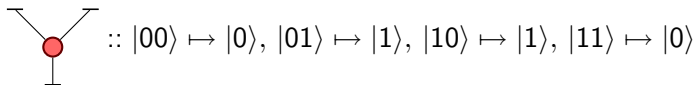
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- ▶ We can depict  $(\delta_X, \epsilon_X)$  with red dots, and note  $\delta_X^\dagger$  is the XOR



# Interactions, and a Rewrite Theory

- ▶ Red and green are both special Frobenius algebras, but they also interact

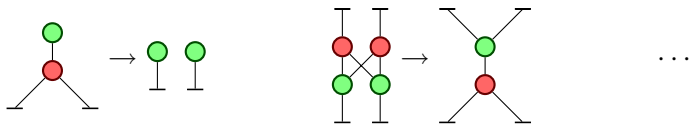
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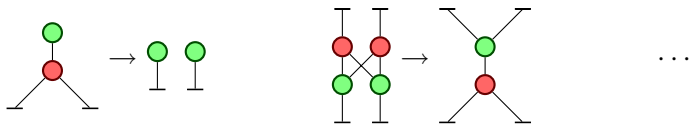
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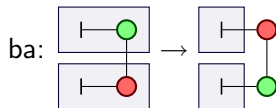
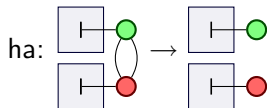
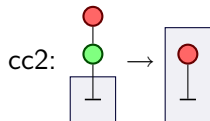
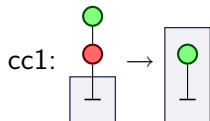
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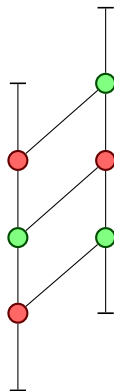
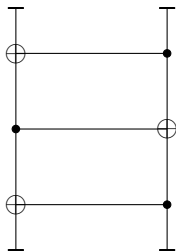
- ▶ We can apply pattern rewrites to pattern graphs, so we get...

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► **Derived** pattern rewrites:



# Doin' the Swap



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- ▶ These yield spider theorems, necessitate pattern rewriting

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- ▶ Quantum information has an abstract graphical language, phrased in your favourite flavour of monoidal category
- ▶ Frobenius algebras pop up everywhere (classical data, entangled states, TQFT's)
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- ▶ Future work: polish the algorithm, more theories, automatic critical pair analysis

# Thanks!

- ▶ This is joint work with
  - ▶ Bob Coecke  
<http://www.comlab.ox.ac.uk/people/bob.coecke/>
  - ▶ Ross Duncan  
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  - ▶ Lucas Dixon  
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- ▶ Check it out at
  - ▶ <http://dream.inf.ed.ac.uk/projects/quantomatic/>